Foundation of proofs

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http://joshua.smcvt.edu/proofs

The need to prove

In Mathematics we prove things

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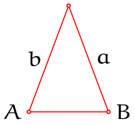

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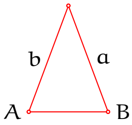

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But what about: 'in a right triangle the square of the length of the hypoteneuse is equal to the sum of the squares of the other two sides'? Is that obvious, or does it require justification?

A characteristic of our subject is that we are completely sure of new results because we show that they follow logically from things we've already established.

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- When decomposed, $18 = 2^1 \cdot 3^2$ has an odd number of prime factors ($1 + 2$ of them), while $24 = 2^3 \cdot 3^1$ has an even number ($3 + 1$ of them). We say that 18 is of *odd* type while 24 is of *even* type.

n	1	2	3	4	5	6	7	8	9
type	even	odd	odd	even	odd	even	odd	odd	even

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Elements of logic

Propositions

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These are not propositions: '$3 + 5$' and 'x is not prime'.

Negation

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So the truth value of ‘not P’ depends only on the truth of P.

Conjunction, disjunction

A proposition consisting of the word and between two sub-propositions is true if the two halves are true.

'$3 + 1 = 4$ and $3 - 1 = 2$' is true

'$3 + 1 = 4$ and $3 - 1 = 1$' is false

'$3 + 1 = 5$ and $3 - 1 = 2$' is false

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A compound proposition constructed with or between two sub-propositions is true if at least one half is true.

'$2 \cdot 2 = 4$ or $2 \cdot 2 \neq 4$' is true

'$2 \cdot 2 = 3$ or $2 \cdot 2 \neq 4$' is false

'$2 \cdot 2 = 4$ or $3 + 1 = 4$' is true

Truth Tables

Write $\neg P$ for 'not P', $P \wedge Q$ for 'P and Q', and $P \vee Q$ for 'P or Q'. We can describe the action of these operators using truth tables.

P	$\neg P$
F	T
T	F

P	Q	$P \wedge Q$	$P \vee Q$
F	F	F	F
F	T	F	T
T	F	F	T
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One advantage of this notation is that it allows formulas more complex than you could say in a natural language. For instance, $(P \vee Q) \wedge \neg(P \wedge Q)$ is hard to express in English.

Another advantage is that a natural language such as English has ambiguities but a formal language does not.

Exclusive or

Disjunction models sentences meaning ‘and/or’ such as ‘sweep the floor or do the laundry’. We would say that someone who has done both has satisfied the admonition.

In contrast, ‘Eat your dinner or no dessert’, and ‘Give me the money or the hostage gets it’, and ‘Live free or die’, all mean one or the other, but not both.

P	Q	P XOR Q
F	F	F
F	T	T
T	F	T
T	T	F

Implication

We model 'if P then Q' this way.

P	Q	$P \to Q$
F	F	T
F	T	T
T	F	F
T	T	T

(We will speak to some subtle aspects of this definition below.) Here, P is the antecedent while Q is the consequent.

Bi-implication

We take 'P if and only if Q' to mean the two have the same values, 'a number n is divisible by 5' if and only if 'the number n ends in 0 or 5'.

P	Q	$P \leftrightarrow Q$
F	F	T
F	T	F
T	F	F
T	T	T

Mathematicians often write 'iff'.

All binary operators

We can list all of the binary logical functions.

P	Q	$P\ \alpha_0\ Q$
F	F	F
F	T	F
T	F	F
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P	Q	$P\ \alpha_1\ Q$
F	F	F
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$\ldots$

P	Q	$P\ \alpha_{15}\ Q$
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These are the unary ones.

P	$\beta_0 P$
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A zero-ary operator is constant so there are two: T and F.

Evaluating complex statements

No matter how intricate the propositional logic sentence, with patience we can calculate how the output truth values depend on the inputs. Here is the work for $(P \to Q) \land (P \to R)$.

P	Q	R	$P \to Q$	$P \to R$	$(P \to Q) \land (P \to R)$
F	F	F	T	T	T
F	F	T	T	T	T
F	T	F	T	T	T
F	T	T	T	T	T
T	F	F	F	F	F
T	F	T	F	T	F
T	T	F	T	F	F
T	T	T	T	T	T

The calculation decomposes the statement into its components $P \to Q$, etc., and then builds the truth table up from the simpler components.

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An important example is that $P \to Q$ and $\neg Q \to \neg P$ are equivalent.

P	Q	$P \to Q$	$\neg Q$	$\neg P$	$\neg Q \to \neg P$
F	F	T	T	T	T
F	T	T	F	T	T
T	F	F	T	F	F
T	T	T	F	F	T

Discussion: our definition of 'implies'

For $P \to Q$ everyone expects when P is true then Q will follow, so that if P is T but Q is F then the statement as a whole is F. What about the other cases?

P	Q	$P \to Q$
F	F	
F	T	
T	F	F
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Standard mathematical practice defines implication so that, for instance, this statement is true for all real numbers:

if x is rational then x^2 is rational

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The intuition is that $P \to Q$ is a promise that if P holds then Q must hold also. If P doesn't hold, that is not a counterexample to the promise. If Q does hold, that is also not a counterexample.

Points about implication

P	Q	$P \to Q$
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- If the antecedent P is false then the statement as a whole is true, said to be vacuously true. If the consequent Q is true then the statement as a whole is true.

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- In particular, our definition does not require that the antecedent P causes, or is in any way connected to, the consequent Q.
- Truth tables show that $P \to Q$ is logically equivalent to $\neg(P \wedge \neg Q)$, to $\neg P \vee Q$, and also to the contrapositive $\neg Q \to \neg P$.

Predicates, Quantifiers; complete statements

Here is a typical mathematical statement (it happens to be false).

$$\text{If } n \text{ is odd then } n \text{ is a perfect square.} \qquad (*)$$

It involves two clauses, 'n is odd' and 'n is square'. For each, the truth value depend on the variable n.

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A predicate is a truth-valued function. An example is the function *Odd* that takes an integer as input and yields either T or F, as in *Odd*$(5) = T$. Another example is *Square*, as in *Square*$(5) = F$, that tells if the input is a perfect square.

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A quantifier delimits for how many values of the variable the clause must be true, in order for the statement as a whole to be true.

Besides 'for all' we will also use 'there exists', denoted $\exists$. The statement $\exists n \in \mathbb{N}\big[\mathit{Odd}(n) \to \mathit{Square}(n)\big]$ is true.

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- There are five different powers n where $2^n - 7$ is a perfect square.

$$\exists n_0, \ldots, n_4 \in \mathbb{N}\big[(n_0 \neq n_1) \wedge (n_0 \neq n_2) \wedge \cdots \wedge (n_3 \neq n_4)$$
$$\wedge\ \exists a_0 \in \mathbb{N}(2^{n_0} - 7 = a_0^2) \wedge \cdots \wedge \exists a_4 \in \mathbb{N}(2^{n_4} - 7 = a_4^2)\big]$$

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- Any two integers have a common multiple.

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- The function $f\colon \mathbb{R} \to \mathbb{R}$ is continuous at $a \in \mathbb{R}$.

$$\forall \varepsilon > 0\ \exists \delta > 0\ \forall x \in \mathbb{R}\ \left[(\,|x-a| < \delta) \to (|f(x) - f(a)| < \varepsilon)\,\right]$$

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is

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which is equivalent to this.

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Thus a person could show that 'every odd number is a perfect square' is false by finding a number that is both odd and not a square.

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